

Understanding and Analyzing Java Performance

Tutorial - MASCOTS 2001

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Outline

- Introduction to Java (10 mts)
- Motivation for studying Java Performance (15 mts)
- Overview of Java Architecture (30 mts)
- Break (5 mts)
- Impact of Java Architecture on Performance (45 mts)
- Analyzing Java Programs (15 mts)



Background



Java : A Brief History

- Appeared in 1993
- Initially developed for
 - Networked, handheld devices
- Coincided with emergence of WWW

 Target market shifted to Internet
- First showcased in the form of applets



...Java- A Brief History

- Server-side Java applets or Servlets introduced as demand increased for dynamic page generation
- Java Beans for reusable software components
- Java Server Pages for decoupling dynamic data from HTML presentation
- "Java 2" Java 1.3

July 27, 2001 HotSpot Compiler, enhancements



Java - Features

- Platform independence
- Security
- Robustness
- Network mobility
- Multi-threaded
- Built in memory management
- Rich API for Internet and Web
 Programming



Why Java ?

- Faster, less troublesome development
- Easy porting to multiple platforms
- Easier software distribution
- Security features
- Rich APIs (Internet, Web,...)

API = Application Programmer's Interface



Where's the catch ?

- Performance !
- Generally true : rich programming features come at the cost of performance. Solutions:
 - Do not use rich environments
 - Understand the environment and do "enlightened" development



Other Disadvantages

- Buggy Virtual Machines
- "Write once Debug Everywhere"
- Platform Independence → "independence" from useful OS features

Bottom Line: Java has become extremely popular, equally among new programmers as well as seasoned C++ gurus.



Motivation



Simple Example: Java vs. C++

class Salutation Java with "Just-in-time" compilation and without

```
private static final String hello = "Hello, world!";
private static final String greeting = "Greetings,
planet!";
```

```
private static final String salutation = "Salutations,
orb!";
```

```
private static int choice;
```

```
public static void main(String[] args) {
```

```
int i;
```

```
for (i = 0; i <= 10000; i ++) {
```

```
choice = (int) (Math.random() * 2.99);
String s = hello;
if (choice == 1) {
    s = greeting;
}
else if (choice == 2) {
    s = salutation;
}
System.out.println(s);
}
```

Java

#include<iostream.h> #include<stdlib.h> #include<string.h>

```
void main () {
```

```
char* hello = "Hello World!";
char* greeting = "Greetings, Planet!";
char* salutation = "Salutation, Orb!";
char* s;
```

```
int choice;
```

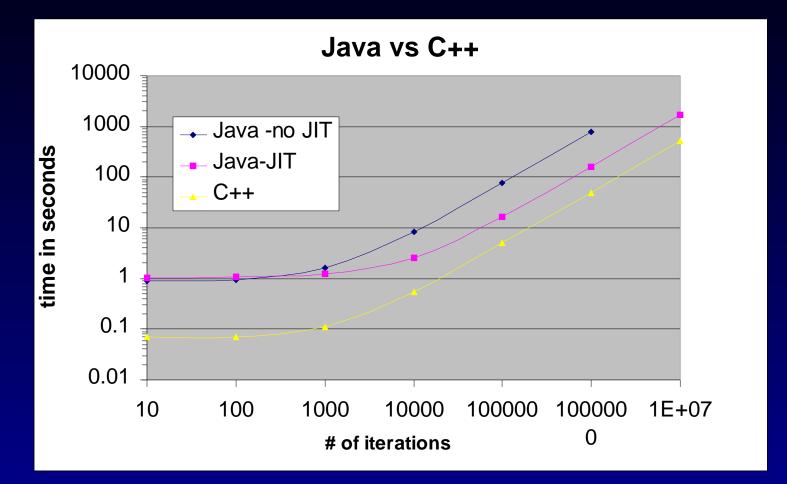
int i;

```
for (i =0; i <= 10000; i ++) {
    choice = ((int) rand()) % 3;
    s = hello;
    if (choice == 1) {
        s = greeting;
    }
    else if (choice == 2) {
        s = salutation;
    }
    cout << s << endl;
}</pre>
```

C++



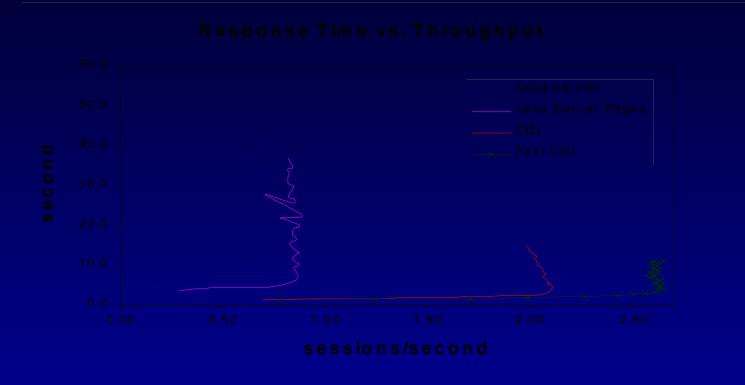
Java vs C++ Simple Example





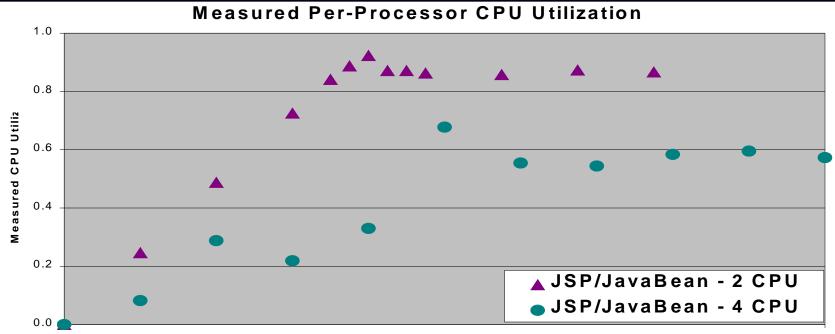
Is Java Slow ?

Realistic Example : A messaging application implemented in Java (servlets and JSP) and C++ (CGI and FastCGI) [5]



Does Java have scalability problems?





Number of Simultaneous Simulated Users

Figure from "Implications of Servlet/Javabean technology on Web server scaling": Cura, Ehrlich, Gotberg, Reeser

Bottleneck prevents use of multiple CPUs efficiently
Thorough analysis pointed to inherent Java bottleneck



Java scalability

- Some history of poor scalability: e.g. Java 1.1.7
 - Article in JavaWorld, August 2000 "Java Threads may not use all your CPUs", P. Killelea.
 - Two programs:one in C, that does an empty loop, same in Java.
 - Run the program as multiple processes on 12-CPU machine scalability of C++ process
 - Run the Java program as multiple threads



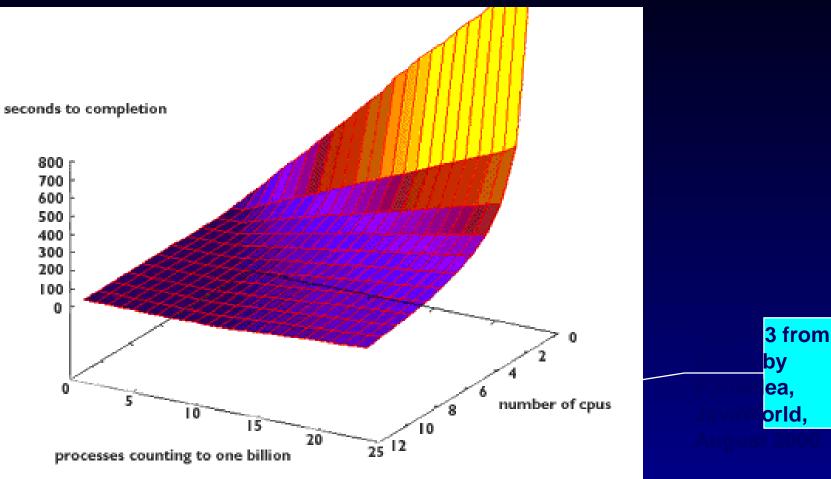
Java Scalability

```
The C program:
main() {
    unsigned long i
    for (i = 0; i < 100000000; i ++);
}</pre>
```



```
The Java program:
class Loop implements Runnable {
public static void main (String[] args) {
  for (int t = 0; t < Integer.parseInt(args[0]); t++)
        new Thread(new Loop()).start();
}
public void run() {
  for (int i = 0; i < 100000000; i ++);</pre>
```

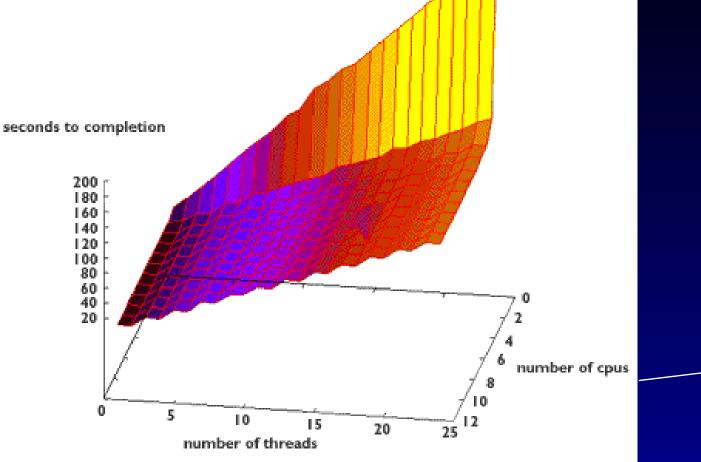
CPU Scalability -C processes







CPU scalability-Java Threads







Initial Conclusion

- Java has performance problems

 Root cause often hard to understand
- But Java has immense technical and business advantages
 - Use of Java for server programs will continue increasing
- Developers and analysts need to educate themselves on Java architecture and performance



Tutorial Goal

- Basic understanding of how Java works
- Identify elements of Java
 architecture that impact performanc
- Intro to issues in performance analysis of Java programs
- Guidelines to improving Java performance (references, papers, etc)



Java Architecture



How Java Works

- 1. Write code in Java : foo1.java, foo2.java
- 2. Compile:
 - javac foo1.java foo2.java
 - (javac is the Java compiler)
 - generates bytecodes in a class file:
 - foo1.class, foo2.class
- 3. Run:

Each application runs inside its own JVM

– java foo1.class

("java" is the JVM: Java virtual machine)

Note: No linked executable



Java Platform Components

- Programming Language
- Class file format
- API
- JVM
- JVM+API = platform for which Java programs are compiled



Programming Language

- Object Oriented
- Robustly checked (type checking, array bounds, memory references...)
- No explicit memory management functions (no free(), destroy())
- Syntactically like C++
- Has a rich class library vectors, hastables, Internet, Web, ...
- Naturally multithreaded



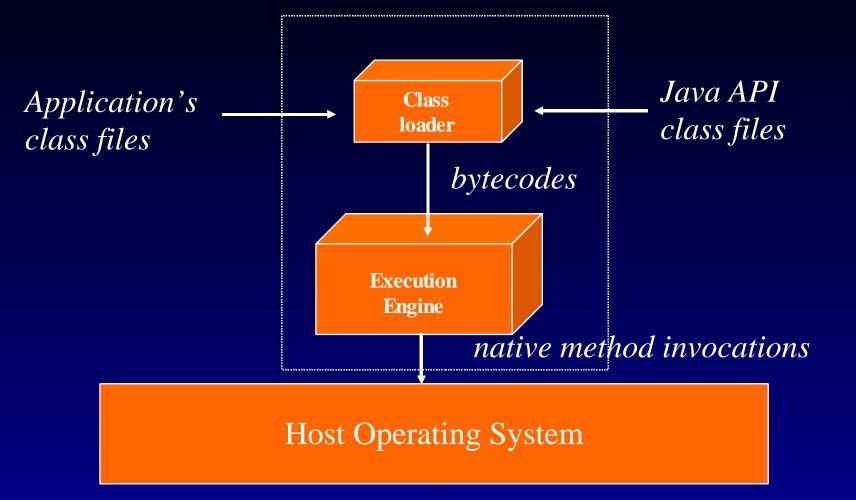
Java Class File

- Binary file format of Java programs
- Completely describes a Java class
- Contains bytecodes the "machine language" for a Java virtual machine
- Designed to be compact

 minimizes network transfer time
- Dynamically Linked
 - can start a Java program without having all classes - good for applets



The Java Virtual Machine



*Figure 1-4, from Venners[1]



JVM (Java Virtual Machine)

- JVM Class loader loads classes from the program and the Java API
- Bytecodes are executed in the execution engine
- Interpreted or
- just-in-time complied : method compiled to native instructions when first compiled, then cached



The Java API

- Set of runtime libraries that provide a standard way to access system resources on a host machine
- JVM+Java API are required components of the Java Platform
- The combination of loaded class files from a program, the Java API and any DLLs constitutes a full program executed by the JVM



Java Under the Hood



Java VM architecture

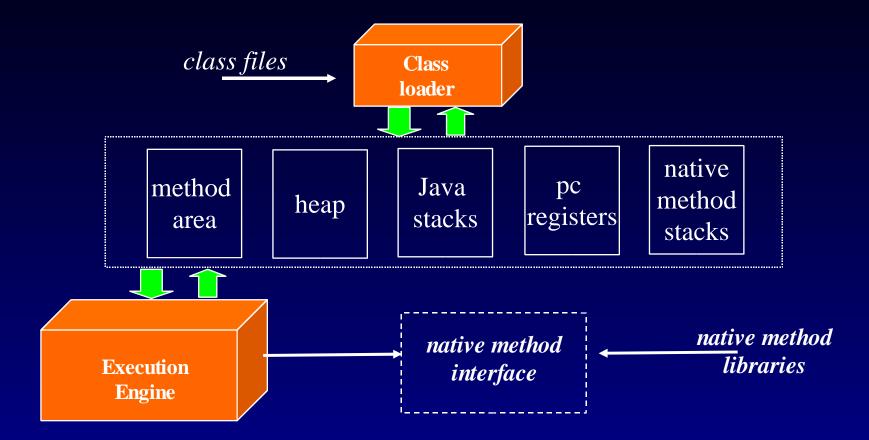
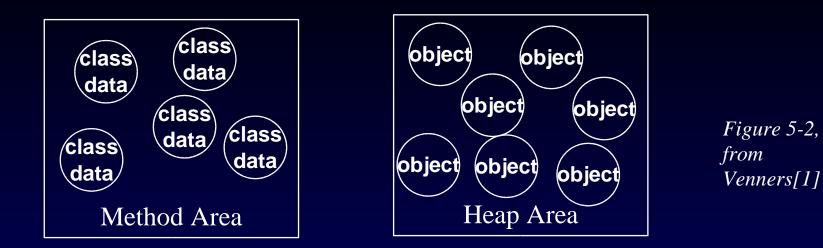


Figure 5-1, from Venners[1]

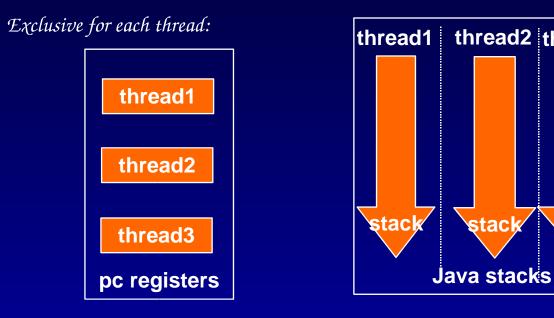
Permormance Analysis JVM: Run-Time Data Areas

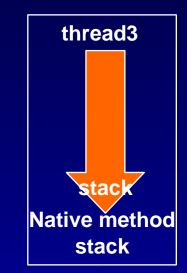
Shared:



thread2 thread3

stack stack



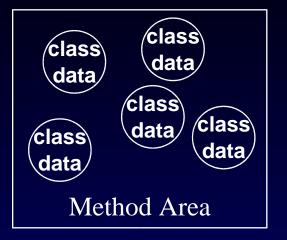


twork Design and

Figure 5-3, from Venners[1] 32



Method Area: • Class loader



•Class loader loads class information in this area

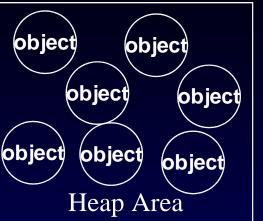
- •All threads share the same method area
- must be thread-safe
 - If one thread is loading a class, the other must wait
- •Method area could be allocated on the heap also
- •Can be garbage collected
 - -Collect unreferenced classes

•Type information:name, superclass name,field info, method info, method bytecodes, a reference to class "Class",...



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The Heap

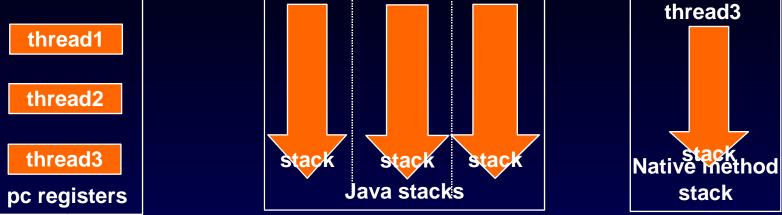


- Area where memory is allocated for objects created during run-time
- Each object has instance data, and pointer to class data in the method area
- Not shared between two applications (each runs inside its own JVM)
- Shared between multiple threads of the same application
 - Access to heap must be thread-safe
 - Access to objects must be thread-safe
- Is managed by JVM using automatic garbage collection (GC)
 - Memory from unreferenced objects is reclaimed
- May have an associated handle pool that points to the actual objects
 - Object reference: Pointer into handle pool

July 27, 200 When objects are moved during GC - update only the handle pool



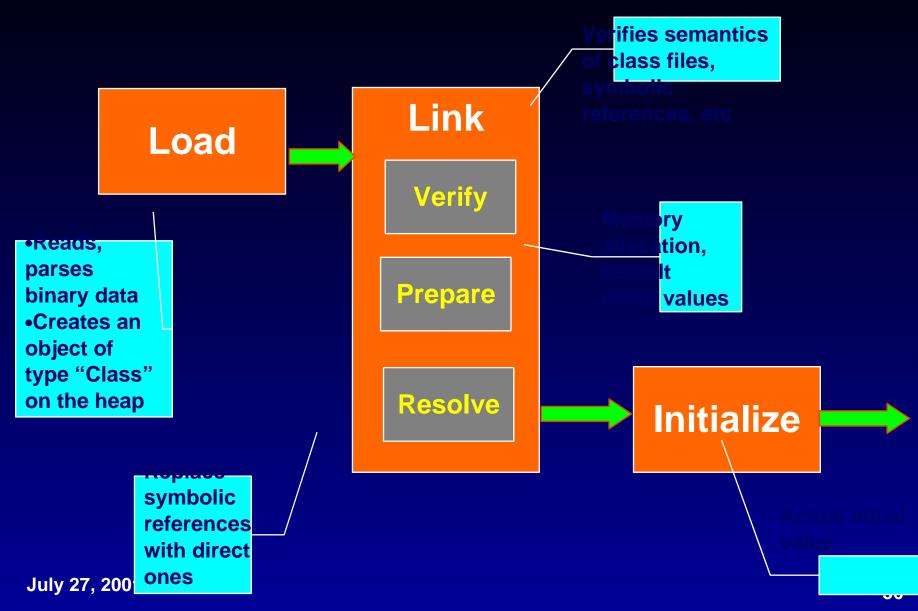
Stacks, PCs



- Each thread has separate stack no danger of access by another thread
 - no danger of access by another thread
- Method calls generate stack frames containing parameters, local variables etc
 - may also be allocated on the heap

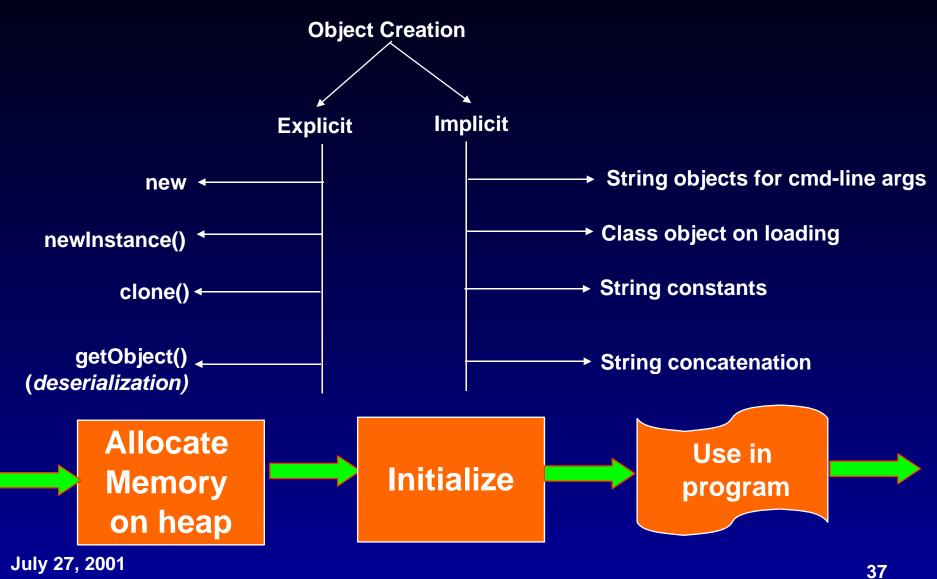


Lifetime of a class





Class Instantiation

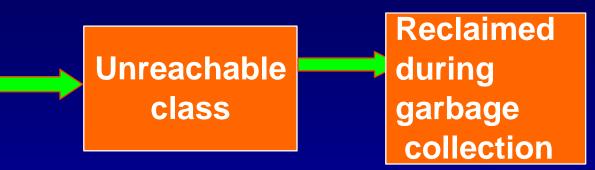




Discarding objects



Discarding classes



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Garbage Collection

- JVM recycles memory used by objects that are no longer referenced
- GC needs to
 - Determine which objects can be freed, free them
 - Take care of heap fragmentation
- Various algorithms for GC, JVM specification doesn't force any one.



Garbage Collection Algorithms

• Tracing Collectors:

 Trace from roots (e.g. local variables, operands) down the reference graph.
 Collect unreachable objects

- Counting Collectors:
 - Maintain reference count for objects.
 - Collect when count goes down to zero.
 - Cannot detect circular references



Garbage Collection -Heap Compaction

- Compacting Collectors:
 - Slide live objects over to occupy free space
- Copying Collectors

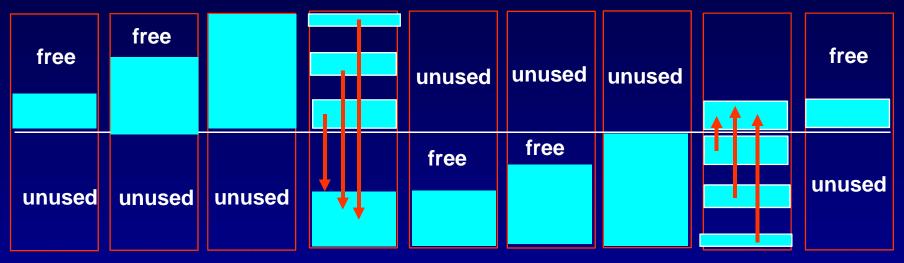


Figure 9-1- from Venners[1]



Garbage Collection -Compaction

- Generational Collectors: Two
 observations:
 - 1. Most objects are short-lived
 - 2. Some objects have long lives
 - Group objects by age or "generations"
 - GC younger generation more frequently
 - Surviving objects move up generations



Synchronization

- Java has a multi-threaded architecture
 - Easy to write code that will not work well with multiple threads
- Use synchronization constructs for
 - Mutual Exclusion: For coherent use of shared data
 - Synchronized statements
 - Synchronized methods
 - Co-operation
 - Working together towards a common goal
 - wait and notify commands



....Synchronization

 Implemented by acquiring locks on objects

```
- Synch statements - lock any object
```

```
class someClass {
    int someVar;
    synchronized(anObject) {
        someVar++;}
```

```
    Synch methods - lock the object on
which the method was called
```

```
class someClass {
    int someVar;
    synchronized void incr {
        someVar++;
    }
}
```

}



Exceptions

- Error handling mechanism
 - programmer can "throw" exception
 - Exception object is created with string comment and stack trace
- Involves object creation, initialization



Security

- Security achieved by:
 - Strict rules about class loading (will prevent loading malicious classes)
 - verification of class files
 - run-time checking by JVM
 - Security manager and the Java API (manages access to resources outside the JVM)



Performance Impact of Java Architecture

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Why is Java slow ?



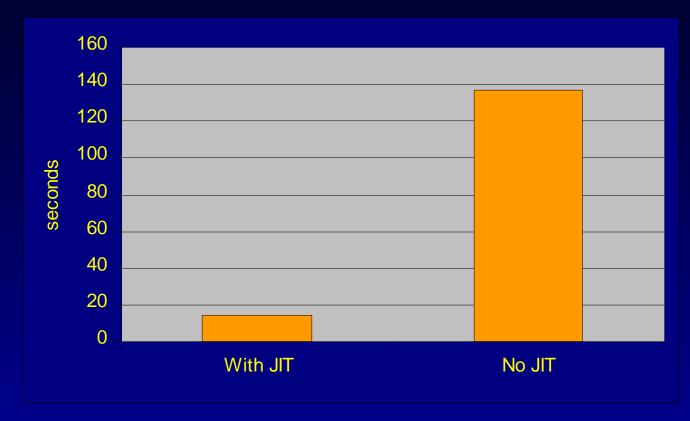
"Obvious" contributors :

- **Bytecode Interpretation (if not jit-ed)**
 - Server-side applications may spend only 10-20 % of time executing Jit-ed code (IBM Systems Journal Paper[3].)
- If jit-ed, compilation cost (one-time), footprint cost
 - OS memory management overhead (paging, scanning etc)



Example

•M/M/1 Queue Simulation: Factor of 10 difference in execution time





More Basic Features Impacting Performance

- Dynamic Linking
- Checking of array bounds on each access
- Checking for null references
- Primitive types are the same- not adjusted to the most efficient type for each platform





Why is Java slow? -Major contributors

- Non-obvious, but deeply impacting performance:
 - Object creation
 - Garbage collection
 - Synchronization
 - API classes too general
 - General-purpose design <u>always</u> implies performance penalty
 - Improper use of Classes and APIs

Performance Impact of Object Creation Object Creation involves:

Allocating memory

 including for superclasses





- Initializing instance variables to Java defaults
- Calling Constructors

including superclass constructors

 Initializing instance variables as programmed
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Performance Impact of Object Creation

 Example 1: Code piece A is 95 % faster than Code piece B

A: boolean bool = a.equalsIgnoreCase(b); **B** :

ucA = a.toUpperCase(); ucB = b.toUpperCase(); boolean bool = ucA.equals(ucB);

• Example 2: Code piece A is 60 % faster than Code piece B

A : Vector v = new Vector(); for (i=0; i<n; i++) { v.clear(); v.addElement... } July 27, 2001

B:

for (i=0; i<n; i++)

{Vector v = new Vector(); v.addElement..}



Object Creation

Two Overheads:

- Creating the object in the heap (previous slide)
- Since the heap is shared by all threads -
 - overhead due to contention for the heap



Object Creation... Scalability

- Concurrency efficiency of object creation across threads
 - Program that creates 500,000 objects, on 6-cpu machine

```
public void run () {
int i;
  myObj obj;
  Thread ct = Thread.currentThread();
  String thrName = ct.getName()+ ":";
  obj = new myObj();
  for (i = 0; i < mt; i++) {
     if (c == 1) obj = new myObj();
  }
}</pre>
```



Object Creation... Scalability

 Time program with varying number of threads- but total # of objects created is always 500,000.

# threads	execution time	"i	"ideal" time	
1	15 s		15 s	
2	10 s		7.5 s	
3	8.5 s		5 s	
4	7.6 s		3.75 s	
5	8.5 s		3 s	
6	8.3 s		2.5 s	



Scalability : Sanity check Concurrency efficiency of cpu-bound program

for (i = 0; i < mt; i++) { for (j = 0; j< 100; j++) f = (i)/ (j+1); }

Timings with varying # of threads (# of loop iterations is constant)

# threads	execution time "ideal" time		deal" time	
1	18.9 s		18.9 s	
2	9.9 s		9.4 s	
3	6.9 s	\approx	6.3 s	
4	5.6 s		4.7 s	
5	4.6 s		3.8 s	
6	4.1 s		3.1 s	



Object Creation

- Observations
 - Has a basic overhead
 - Programs doing lot of object creation (explicit/implicit) will have unexpected scalability problems
 - Each created object adds to garbage collection overhead
 - must be traversed
 - must be collected, when unreferenced.
 - Having many short-lived objects can be a performance bottleneck



- Garbage collection adds a run-time
 overhead
 - In older JVMs GC could stop all processing
 - GC could result in user perceivable delays
 - Delays could be 5-10 seconds for large heaps (100-500 MB)[3]



- Newer JDKs have improved algorithms
 - Sun JDK 1.3 has
 - Generational garbage collection
 - Train algorithm for the old generation subheap
 - Overhead is now smaller
 - e.g. Queue simulation example : 53 ms out of 13 s running time. Heap size b/w 160KB and 2MB
 - Is larger if heap is large



- Garbage collection can be timed (java -verbosegc)
- Test GC in a program in which number of objects, and heap size keep increasing

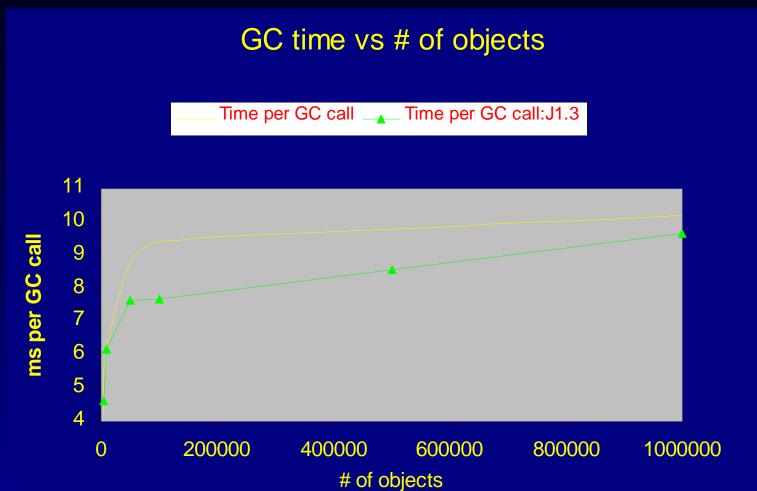
class create implements Runnable {
 static int m, c, mt;
 public void run () {

int i; myObj obj[]= new myObj[1000000]; Thread ct = Thread.currentThread(); String thrName = ct.getName()+ '':'';

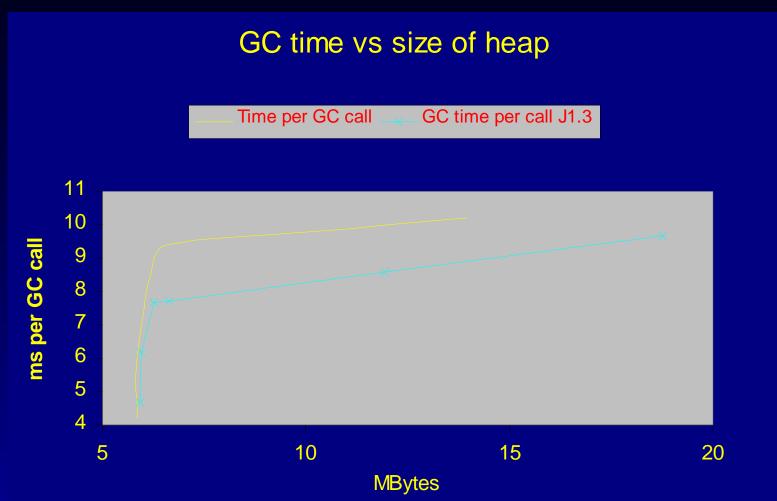
```
long st = System.currentTimeMillis();
for (i = 0; i < mt; i++) {
    if (c == 1) obj[i] = new myObj();
    //System.out.println(thrName+obj);
    }
    long diff = System.currentTimeMillis() -st;
    System.out.println('Time: ''+ diff);
```

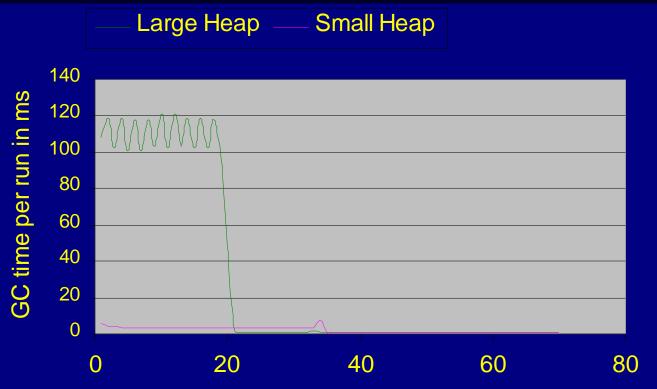
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Test queue simulation program, after allocating a large array of objects in the beginning, and then running the simulation as usual.

rk Design and nance Analysis

GC iteration

•Looks like GC learns about the long-lived object and does not include that in later GC?

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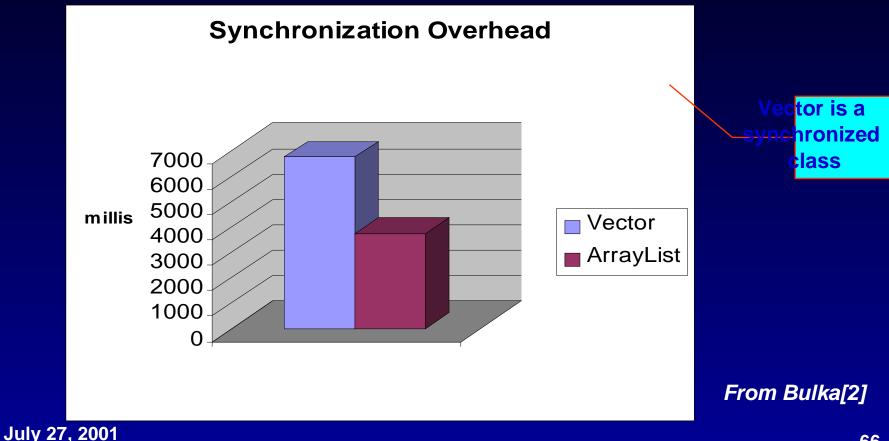
- "Obvious" :
 - In a multithreaded application, synchronized pieces will be the bottlenecks (Javaindependent issue)
- Non-obvious (Java-isms):
 - Big synchronization overhead
 - Java API classes may have synchronized methods - a big overhead in cases where synchronization is not necessary (access only by one thread)

Implicitly shared objects internal to the JVM e.g. heap. Access will be synchronized

Performance Impact enjormance Analysis Something. of Synchronization

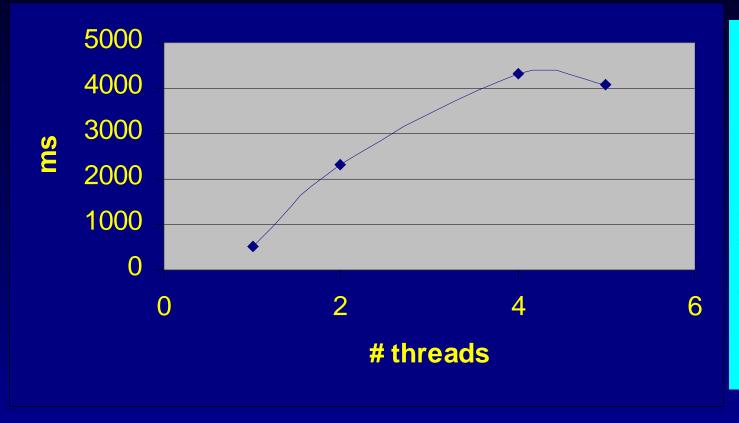
• Example: Vector vs ArrayList (example)

creates vector/array list, adds elements, then accesses them)



Design and

Performance Impact of Synchronization Contention for synchronized code:



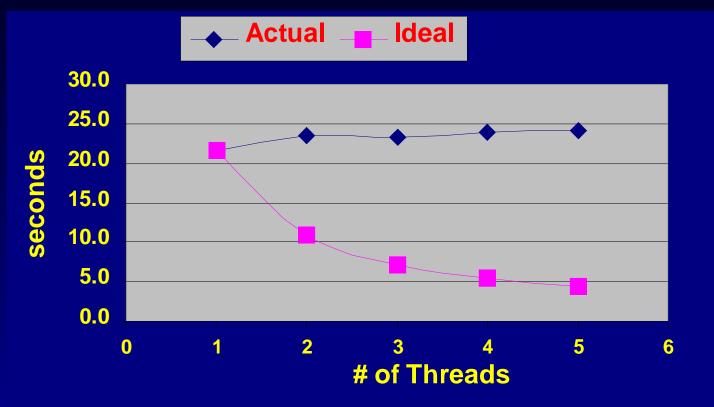
Example Bulka[2]: increase a counter using synchronized method. Use increasing # of threads to do the same amount of total work. Results from 6-cpu machine.

k Design and

Performance Analysis

Contuines

Performance Impact of Synchronization Implicitly synchronized code:



Object creation example, with printing inside the loop (System.out. Println - not an explicitly synchronized function in Java. Access possibly serialized by OS)





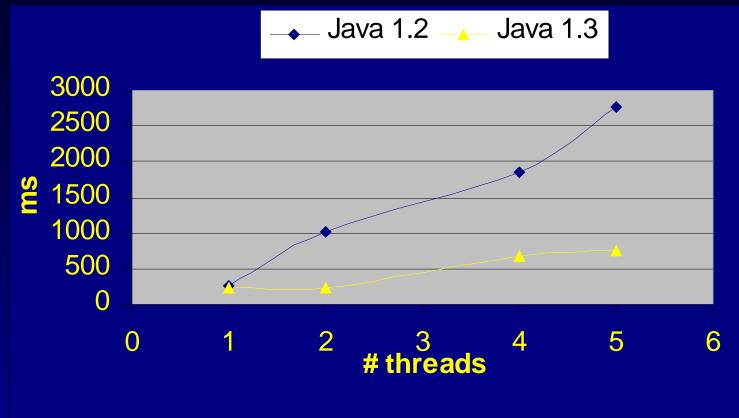
Example: Multiple threads increment a shared variable by calling Math.random()

Run this program with increasing number of threads, keeping the total number of iterations the same - on 6-CPU machine class WorkerThread extends Thread {
 private int iter;
 private int tid;
 private static double num;
 public WorkerThread (int
 iterationCount, int id) {
 this.iter = iterationCount;
 this.tid = id;
 }
}

```
public void run() {
  for (int i = 0; i < iter; i++) {
    num += Math.random();</pre>
```



 Example of multiple threads calling Math.random() - a synchronized method





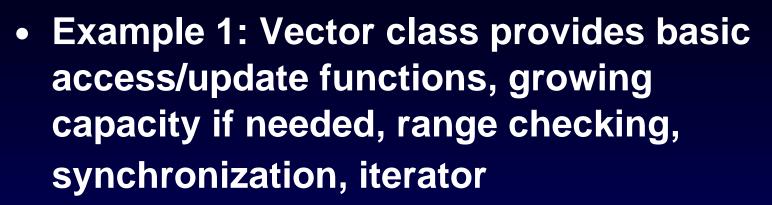
- Object creation can be viewed as a special case of access to synchronized data structures and methods
- We saw similar effects there

General-Purpose API classes



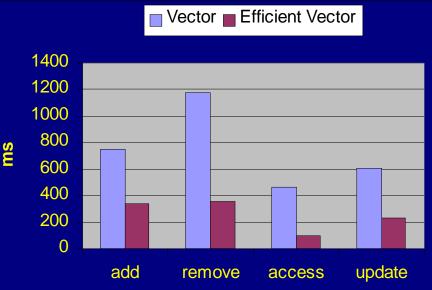
- Generally true: When a class/API provides maximum flexibility and features, there will be an associated performance cost.
 Examples:
 - Vector Class
 - Some applications may need their own efficient vector implementation
 - Date
 - Using native Date functions thru JNI might prove better performing

General-Purpose API classes



Example from Bulka[2]:

Speed up due to a "light" implementation of Vector class, offering few features.



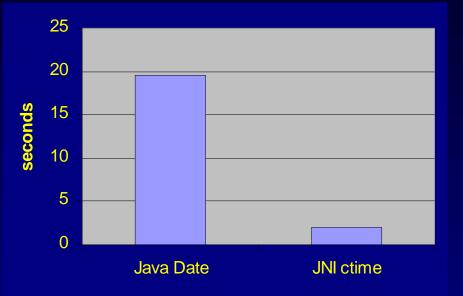
k Design and ance Analysis



Performance Impact of "Heavy" API classes

 Date is a computationally expensive class

Example from Bulka[2]:



Speed up due to a use of native call instead of the Java Date class



Java Memory Issues

- Contributors to memory usage in Java:
 - Objects
 - Classes
 - Bytecode
 - JIT compiled code
 - Constant pool entries
 - Data structures representing methods and fields
 - Threads
 - Native data structures
 - e.g. OS-specific structures
- Too much memory usage will result in OS virtual memory overheads - and possible slow down in garbage collection



Java Memory Issues

- No method for calculating object size
 - Methods returning total memory and free memory of heap
 - Object size can be estimated indirectly using garbage collection, and heap memory methods
- Class loading can be tracked with java -verbose: lists all the classes being loaded



Key Recommendations

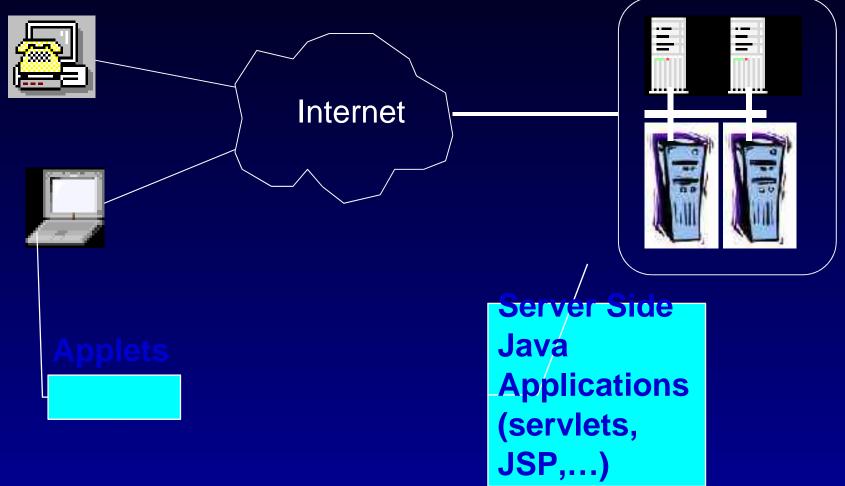
- Limit object creation (various techniques...)
- Do not use synchronized API classes if not needed
- Rewrite "heavy" API classes, if light ones are needed
- Apply various optimizations (books, papers).



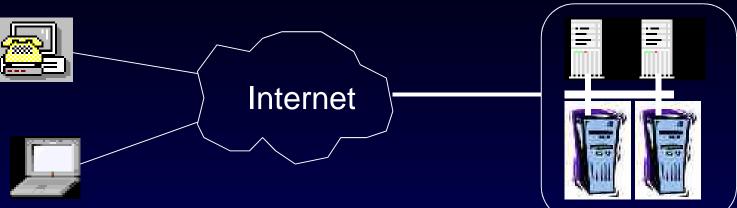
Performance/Capacity Analysis of Java Applications



Two kinds of Java apps



Applet performance Issues



- Download time
 - downloads can be sped up using *jar* files instead of individual class files
- Dynamically linked classes that are downloaded when needed (will affect user response time on first use)
- Needs to be fast (usually used as a GUI)
- Usually no thread contention issues





Capacity Analysis for Server Applications

- Typical industry problem:
 - Given a Java server application, size the server center to support volume of *N* requests per second.
 - Available data: measurement data from load testing at smaller volume and on systems smaller than "production" systems.



Issues in Java App capacity analysis

- Bottleneck capacity may not be that of a hardware resource
- Bottleneck may be
 - a piece of synchronized code
 - object creation, if a large number of objects are being created.
 - garbage collection, if large number of short-lived objects.
 - I/O (poorly coded)

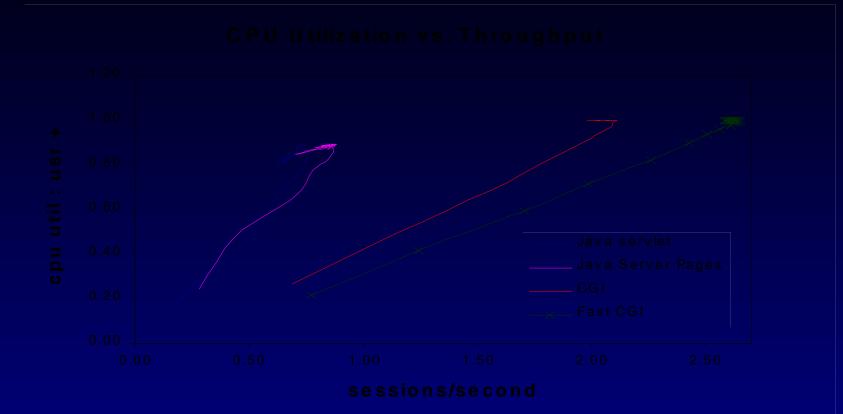


Issues in Java App capacity analysis

- Possibly no capacity increase with additional processors (threads)
 - CPU may not be the bottleneck
- Speed up due to more memory
 _ Configure larger heap size
- Speed up with more servers
- CPU time per transaction may increase going from small to large number of users



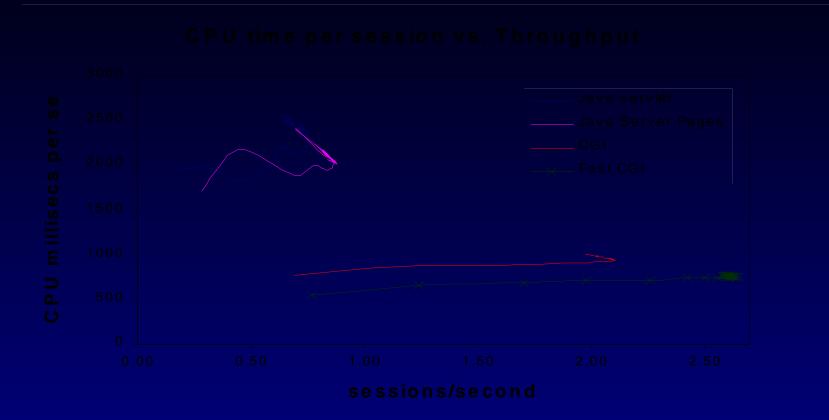
Messaging Example



From Hansen, Mainkar, Reeser, 2001 [6]



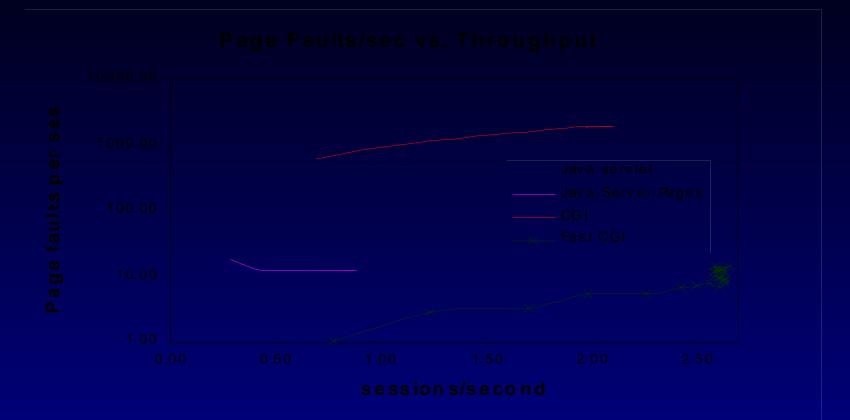
Messaging Example



From Hansen, Mainkar, Reeser, 2001 [6]



Messaging Example



From Hansen, Mainkar, Reeser, 2001 [6]



Delay Analysis

- Apart from hardware resources, Java's software resources should also be analyzed as queues -
 - should take into account synchronized portion of code, and contention for it in a delay model.
- Should take into account garbage collection - service time in queues may be load-dependent



Previous Work

- Reeser[5] modelled a Java application with software "code lock" as a separate queue
 - "Abstract" bottleneck, paper does not say which particular Java resource was the bottleneck
 - Model fits well



Reeser model example

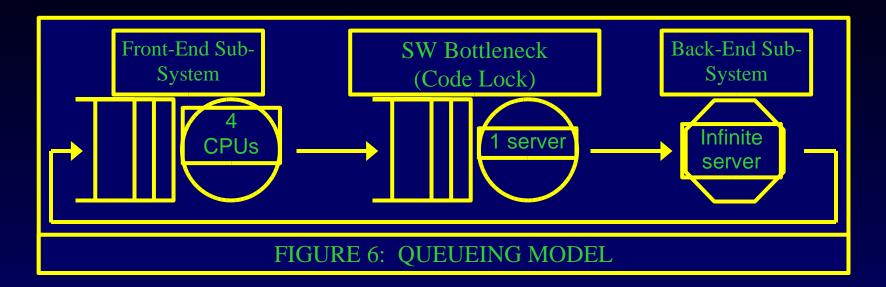


Figure 6 from "Using Stress Test Results to Drive Performance Modeling: A Case Study in "Gray-Box" Vendor Analysis", ITC-16, Brazil, 2001.



Reeser model example

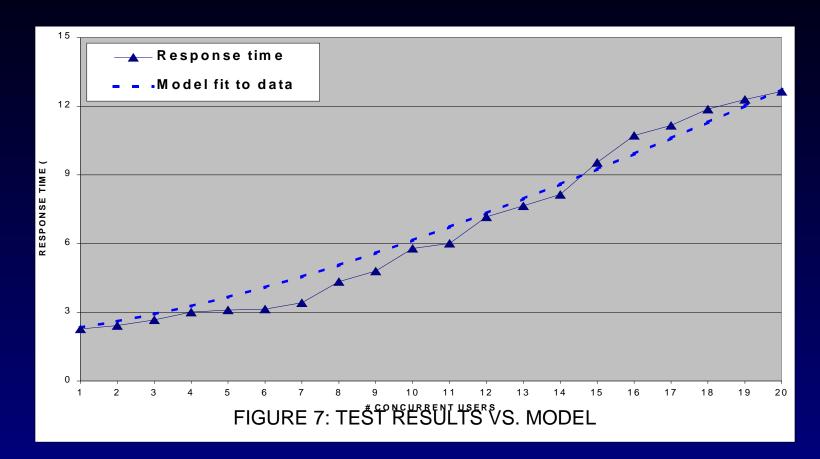


Figure 7 from "Using Stress Test Results to Drive Performance Modeling: A Case Study in "Gray-Box" Vendor Analysis", ITC-16, Brazil, 2001.



Profiling Tools

- Java VM comes with a profiler
 - Can report times spent in method calls, heap data etc.
 - Hard to read and understand
- Commercial Profilers
 - Jprobe, Optimizelt
- Useful to developers to really tune
 their code
- Useful to analysts for understanding July 27, 2007 C time and other bottlenecks



Future Directions

 Better models and techniques to analyze and predict capacity and performance of Java applications



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- 6. T. Hansen, V. Mainkar, P. Reeser, "Performance Comparison of Dynamic Web Platforms", SPECTS 2001.